A Distributed Transaction and Accounting Model for Digital Ecosystem Composed Services

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Abstract—This paper addresses two known issues for dynamically composed services in digital ecosystems. The first issue is that of efficient distributed transaction management. The conventional view of transactions is unsuitable as the local autonomy of the participants is vital for the involvement of SMEs. The second issue is that of charging for such distributed transactions, where there will often be dynamically created services whose composition is not known in advance and might involve parts of different transactions. The paper provides solutions for both of these issues, which can be combined to provide for a unified approach to transaction management and accounting of dynamically composed services in digital ecosystems.

Index Terms—Distributed Services, Coordination, Transactions, Accounting, Digital Ecosystem.

I. INTRODUCTION

Many in the business and research communities are pursuing the vision of digital business ecosystems – distributed software environments through which organisations can seamlessly access customised, potentially disposable, services to aid them carry out a myriad of tasks. These services can be either pure software applications, or “real world” services represented by a software wrapper supporting automated transaction processing. Full realisation of this vision requires deployment of facilities for the dynamic discovery, composition, interoperation and execution monitoring of a potentially huge number of available services. The DBE integrated project \(^1\) focuses on the development of an open-source distributed environment (the DBE) that can support the spontaneous creation of applications through the composition of (not necessarily open-source) software services and components. In doing so the project is adopting the business modelling approach described above, but complementing it with the adoption of evolutionary algorithms inspired by biological processes, that provide bottom-up incremental improvement of business models through run-time feedback on service performance. The DBE is being targeted primarily towards SMEs, will be able to concatenate their offered services within service chains formulated on a pan-European basis. By offering access to a large pool of service providers and consumers, and itself providing advanced recommendation systems and evolutionary algorithms, the DBE will support continued global optimisation of service chains, benefiting all actors, in particular SMEs \(^2\).

However, the adoption of such a collaborative environment by SMEs is largely dependent on whether we are able to guarantee consistency and preserve local autonomy. Current frameworks \(^6,7\) tend to focus on providing consistency but, at the transaction level, this comes at the cost of violating the local autonomy of the participants.

Businesses utilising DBE services will need to be charged and billed in accordance with their service usage patterns. In any business environment charging processes are of crucial importance, thus the success of the DBE, or indeed any environment supporting dynamic service composition, will be contingent on the ability of providers to charge and collect fees for service usage. Given that there will be no a priori knowledge of the structure of, or business model associated with, dynamically composed services, the design and deployment of a sufficiently flexible accounting system is a challenging task.

Both the transaction management and accounting solutions presented here are being made available through open source projects providing SMEs with affordable solutions to engage in eCommerce through digital ecosystems platforms.

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II. DISTRIBUTED TRANSACTION MANAGEMENT

The long-term nature of business transactions frames the concept of a transaction in Digital Business Ecosystems and makes defining a consistent transaction model even more challenging. The conventional view of a transaction \(^5\), based on the ACID (Atomicity, Consistency, Isolation, Durability) properties, cannot capture the primary requirements of DEs. From a business point of view, most usage scenarios in Digital Ecosystems involve long-term transactions and thus Atomicity is an unacceptable constraint. From a distributed transactions point of view, Isolation can lead to significant degradation of performance in the services offered (critical data is locked until a transaction completes) or to increased probability of deadlock (as services may be locked into composite transactions that do not terminate).

Transaction models for web services include WS-BusinessActivity \(^6\) and BTP \(^7\) among others. They are primarily concerned with consistency but in guaranteeing it, the underlying coordination framework requires access to the internal build-up of the communicating parties. The Coordinator and Participant roles are tightly-coupled and if a
fault occurs while the transaction is finalising (transition to Fault state while in Close state) it can only be dealt with at the participant’s platform. This not only implies that participating SMEs cannot have exclusive control of their local design but also that the underlying services are no longer stateless. Such presumptions are against the primary requirements of SOA [9] and hinder the development of a Digital Ecosystem for SMEs.

III. ACCOUNTING AND CHARGING FOR NETWORKED SERVICES

Accounting involves the collection and analysis of service and resource usage metrics for purposes such as billing, capacity and trend analysis, cost allocation and auditing. It requires that service consumption be measured, rated, and that resultant charging information be communicated between appropriate business entities. As shown in Fig. 1, accounting systems for networked services incorporate sub-systems for metering, mediation, rating and billing.

![Fig. 1 Accounting System for Networked Services](image)

Accounting for services, including composite services, whose characteristics are known in advance is a mature area. For example, in the telecommunications domain there is wide deployment of complex mediation, rating and billing systems which support sophisticated usage- and content-based charging models for pre-paid and post-paid, private and corporate customers. In such systems the components involved in the accounting process must be manually pre-configured to account for specific services at the time those services are initially deployed. This will not be the case for composed services dynamically created and executed within a short time span – business logic for accounting must be automatically configured when service compositions are initially constructed, or subsequently modified.

Despite its importance, accounting for composed services has received little attention in the published literature. Bhushan et al. [3] propose a system architecture that supports the requirement for service providers to cooperate in the provision of composed services in a federated manner and share the generated revenue. However, they addresses only statically composed services – accounting components would still have to be pre-configured with the relevant accounting logic. Agarwal et al. [4] propose a method for metering and accounting for composite e-services that is not dependent on a-priori knowledge of the service composition. However, their approach supports only two specific service pricing models (flat rate per amount of resource used and flat rate per transaction). More significantly, the charge for an invocation of a composed service will always be the summation of the charges associated with standalone invocations of the constituent services; this will not always reflect the potentially complex business relationships between the service providers.

IV. DISTRIBUTED TRANSACTION MODEL

Our primary concern is with the support for long-term business transactions involving open communities of SMEs. Hence, service composition in this context is multidimensional and different forms of composition are needed to satisfy evolving business requirements. In contrast with the conventional view of transactions (data centric), web service compositions deal with at least three aspects: order, dependency, and alternative service execution [11].

In our approach, a multi-service transaction is represented by a tree structure (see Fig. 3). Each node of the tree is either a coordinator (can be considered a composed service) or a basic service. The tree describes the coordination of the involved services and we may thus refer to the root of the tree as the master composed service. Basic services appear only as leaves of the tree. The resulting directed graph indicates the sub-transactions involved in a transaction and the corresponding hierarchical collection of basic and composed services. The coordinator nodes determine the ordering of execution and infer dependencies within the associated service hierarchy. We will describe how such issues, often collectively wrapped up in the term choreography [10], are addressed within our approach in the sequel.

First, we consider five different coordinator types, drawing upon [11], that allow for various forms of service composition to be expressed in our model.

1. **Sequential coordinator**: the services are invoked sequentially and the execution of a service is dependent on the previous one. This coordinator can handle sequential process-oriented service composition with provision for both Sequential with Commit Dependency (SCD) and Sequential with Data Dependency (SDD).

2. **Parallel coordinator**: the services can be executed in parallel. This coordinator handles parallel process-oriented service composition covering Parallel with Commit Dependency (PCD), Parallel with Data Dependency (PDD) and Parallel without Dependency (PND).

3. **Sequential Alternative** coordinator: the services will be attempted in succession until one produces the desired outcome, as specified by some criterion (e.g., cost, time)

4. **Parallel alternative** coordinator: alternative services are executed in parallel and once a service produces the desired outcome, the rest are aborted.

5. **Data-oriented** coordinator: this coordinator handles data-oriented service composition and specifically deals with released data items within a transaction (between its sub-transactions) or partial results released between different transactions.

6. **Delegation**: this coordinator allows the whole transaction or a sub-transaction to be delegated to another platform.

The first four coordinator types are rather self-explanatory. In long-lived transactions, partial results need to be shared between transactions before their termination (commitment). This is the purpose of the data-oriented coordinator. The delegation coordinator provides a means of overcoming traffic bottlenecks or low bandwidth connections or (lack of) processing power. In this paper we are
mostly interested in a charging scheme for distributed multi-service transactions and thus these coordinator types will not be covered in greater detail.

Consider the simple transaction tree depicted in Fig. 2 where we have adopted the notation of [8]. It comprises five basic services whose order of execution is determined by the three coordinators.

![Fig. 2 Multi-service transaction in a tree structure](image)

The “master composed service” is a sequential alternative coordinator defining two alternative execution scenarios: (i) service s1 followed by s2, and (ii) service s1 followed by s2, followed by s3, s4, and s5 in parallel. The second scenario of composed services will only be executed if the outcome of the first does not meet some preset condition.

The tree structure representation of a transaction allows us to exemplify the local coordination of the corresponding (compositions of) services that is required in performing the transaction in question.

To accommodate distributed long running transactions that involve composed services, we need to relax the ACID properties, particularly Atomicity and Isolation without compromising Consistency. For this purpose, we need to consider some additional structure that guarantees the consistency of the transaction model within a highly dynamic and purely distributed environment of a Digital Ecosystem.

At the same time, we are considering SOC [9] as the primary computing paradigm for DBE. This entails that the model should defer from any tight-coupling between initiator and coordinator or between initiator and participant, as is the case with WS-BusinessActivity [6]. We therefore keep state information at the deployment level and abstain from interfering with service execution as such as we wish not to break the local autonomy of the realisation platform.

Next, we describe two graphs that capture the dependencies between sub-transactions (basic and/or composite services) of a single transaction or belonging to different transactions. Keeping track of such dependencies is essential if the underlying transaction model is to provide capabilities for charging for services, including composed services, and reverse action (when a sub-transaction fails or is aborted).

The **Internal Dependency Graph (IDG)** is a directed graph of arcs and nodes, which keeps logs of value dependencies within a transaction tree. Each node represents a coordinator or service (sub-transaction) and the direction of the arc between nodes indicates a dependency of one node on another.

In the transaction tree of Fig. 2, for example, s2 and s1 are children of a sequential coordinator and hence s2 is dependent on the results released by s1. This means that service s2 cannot be invoked before s1 has. It also has as a consequence that if s1 is aborted, then s2 must also be aborted. This dependency between s2 and s1 is shown in the IDG of Fig.3(i).

The services s3, s4, and s5 are children of a parallel coordinator. This implies that when value dependencies exist if one of the services is aborted then the rest of the services must also be aborted. This is shown in the corresponding IDG given in Fig. 3(ii).

![Fig. 3 Internal Dependency Graph](image)

We have seen that in a distributed multi-service transaction, dependencies may exist not only between services of a transaction but also between services of different transactions. For instance, consider the case of (compensatable) sub-transactions that release partial results in a conditional commit state [8]. This incurs a dependency between the corresponding service executions.

To capture such dependencies, we use another graph, called the **External Dependency Graph (EDG)**. This graph keeps track of dependencies between (services or coordinators) of different transactions. The log structure it provides can also be used in recovery routines for running a compensating procedure, upon failure.

Fig. 4 shows (part of) the EDG for the transaction T1 (of Fig. 2) and transaction T2. It indicates that the service s8 of T2 uses results released by service s5 of T1, and is thus dependent on s5.

![Fig. 4 External Dependency Graph](image)

The IDG and EDG provide a means of recording important system logs which can be stored locally, on the corresponding local coordinator, but their effect is both local, in terms of local faults, forward recovery and contingency plans, and global, in terms of abortion, restarting, recalculating, and alternative execution. Both graphs are used in the rating process of accounting, discussed in §V, for determining the final charges of composed services of a transaction (IDG) and across transactions (EDG).

### V. TWO PHASE RATING PROCESS

We return now to the rating algorithm. To meet the requirements outlined in §III we propose that DBE rating engines employ a two phase rating process. In phase 1 of
this process, services comprising a composed service are rated as if they are being executed as “standalone” services, leading to the generation of interim charges for each of these services. In phase 2, the interim charges are modified in accordance with provider specified rules specifying how charges are to be modified if that provider’s service(s) are used with other specified services in the context of a composed service. To facilitate the two phase process we utilise two-part charging schemes: part 1 dictates how charges are to be calculated when the service is invoked in isolation, and part 2 dictates any modifications to these charges if other specified services or service providers are present in the current transaction.

For the purpose of rating, we view composed services as hierarchical collections of atomic and composed services, in the same way that the transaction model is illustrated above in Fig. 2. We refer to the service at the top level of the hierarchy as the “master composed service”. Services at the bottom of the hierarchy must all be atomic (basic) services (from the rating engine’s perspective). All DBE services have associated with them a providerID, which uniquely identifies their provider within the DBE; they have a serviceID, which uniquely identifies the service amongst the set of services offered by their provider; and an instanceID which distinguishes between multiple instances of the same service type (for example, serving different geographical markets). Therefore, the tuple of (providerID, serviceID, instanceID) uniquely identifies a service instance across the entire DBE.

When a Transaction Workflow Manager invokes a master composed service it assigns a transactionID that uniquely identifies this invocation of the master composed service during the timeframe between initial invocation and the completion of all processing associated with that invocation. The Transaction Workflow Manager then provides all services comprising the master composed service with this transactionID. All metering records relating to these services will contain this transactionID, as well as the providerID, serviceID, and instanceID. In addition, the Transaction Workflow Manager indicates to the rating engine that an invocation of the master composed service with the specified transactionID is commencing, and provides it with details of the associated service hierarchy. The presence of the transactionID in all metering records provides a means for the accounting system to identify a service’s execution context.

As the services are executed, metering records relating to these records are generated and transferred to the rating engine; these records are rated using part 1 of the charging scheme associated with that service. This is phase 1 of the rating process; it results in the generation of interim charge records for the invocations of the associated service instances.

In the phase 2 of the process the rating engine modifies the interim charge records generated in phase 1 as dictated by the part 2s of the relevant charging schemes. These elements of the charging schemes capture relationships between services or between service providers and provide the rating engine with appropriate discounts (or tariffs) to be applied when such services have transaction dependency relationships as indicated by the IDGs and EDGs provided by the transaction workflow manager.

VI. SUMMARY

Within a Digital Business Ecosystem a number of long running transactions take place, involving various service compositions, and there is an increased likelihood that some sub-transaction is aborted. Preliminary analysis of our transaction model shows that it is possible to provide a compensating mechanism that warranties consistency. The log structures of the IDG and EDG capture the dependencies due to released results within a transaction and partial results between transactions, respectively. To ensure consistency at all times, such dependencies need to be taken into account so that all dependent subtransactions are also aborted. At the same time, it is also possible to address omitted results (non-dependent subtransactions that may have provided valuable results and need not be aborted) through a forward recovery routine, again based on the IDG and EDG.

Additionally, work is in progress on a formal behavioural description of the proposed transaction model, which can be subsequently used for the rating of the corresponding service compositions. By verifying the behaviour exhibited by the underlying service compositions, the testing and deployment of the transactions and their accounting scheme can be eased. The formal representation of a transaction also allows to identify the possible sequences, if any, of reverse actions for compensation and forward recovery.